Recent advances in the Linux kernel resource management

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Agenda

- Resources to account and control
- Some background on containers
- Existing functionality, shortcomings
- Control Groups a.k.a. cgroups
- Memory Controller
- Future work





Why?

- All resources are finite
- Multiple tasks and users
- Need usage statistics / bookkeeping
- Need Denial of Service protection
- Need Quality of Service level (not only limits but guarantees)

What?

- CPU
- Memory (RAM)
- Swap
- Disk space
- Disk I/O
- Network

Resources: CPU

CPU is given to tasks in time slices

- CPU shares/weights
- CPU limits
- for SMP: CPU affinity

Resources: Memory & swap

- User memory
 - Virtual and physical (RSS) memory
 - Dirty page cache
- Kernel memory
 - Various objects, different allocators
 - Special case: network buffers
- Swap space

Resources: disk

- Disk space
- Disk I/O bandwidth
 - read/write
 - mmap()
 - swapin/swapout
 - Problem: most of I/O is async

Resources: networking

- Network bandwidth: solved by tc
- Traffic Control:
 - Shaping
 - Scheduling
 - Policies
 - Dropping

Containers



What are containers?

- Multiple isolated userspace instances
- Running on top of a single kernel
- Like VMs but very lightweight, native performance, low overhead

Containers Implementations

• OpenVZ



- Parallels Virtuozzo Containers || Parallels[®]
- FreeBSD jails
- Linux-VServer VServer
- Solaris 10 Containers/Zones
- IBM AIX6 WPARs (Workload Partitions)







Containers cont'd

- Multiple containers should peacefully co-exist, need DoS protection
- From the resource management point of view, containers are just groups of processes.

Existing mechanisms

Disk Quota

- Per mount point disk quota for users and groups
- Soft limits, hard limits, grace periods
- Can see the current usage
- Can be inc'd/dec'd on-the-fly
- Applications are expecting disk space shortage (or at least should be)

CPU

- Per-process nice value which can be changed on-the-fly (nice, renice)
- Real-time priority queue
- Hard CPU time limit (ulimit -c)

ulimit

- setrlimit()/getrlimit() syscalls
- Controls 16 different resources:

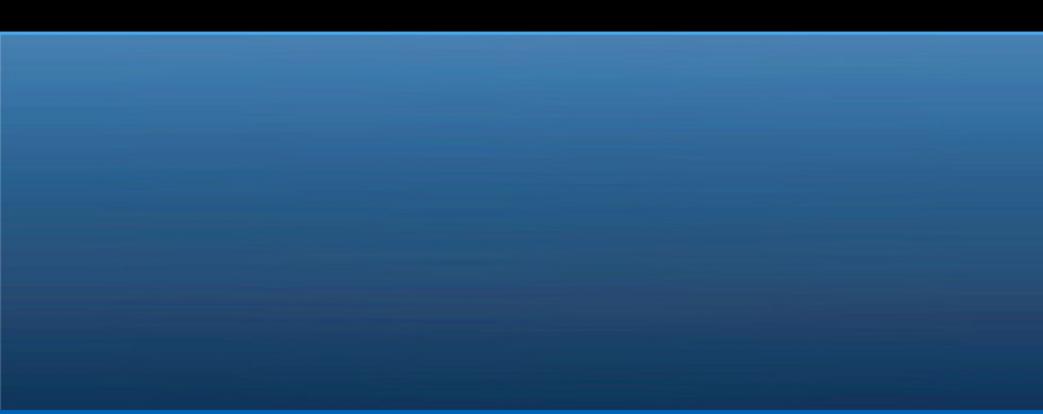
core file size, data seg size, scheduling priority, file size, pending signals, max locked memory, max memory size, number of open files, pipe size, POSIX message queues, real-time priority, stack size, cpu time, max user processes, virtual memory, file locks

Soft limits and hard limits

ulimit: problems

- Not all resources are covered
- Ulimits set in the current context
 - the only good place to set is login
 - some can only be decreased run-time
- All limits are per-process
 - only NPROC is per-UID
- Current usage values are unknown
- Memory limits are mostly ignored

Control Groups



Control Groups

- A generic mechanism for grouping tasks into hierarchical groups
- Multiple resource controllers
- Possible to have different groups for different controllers
- Managed via cgroup filesystem

Control Groups: interface

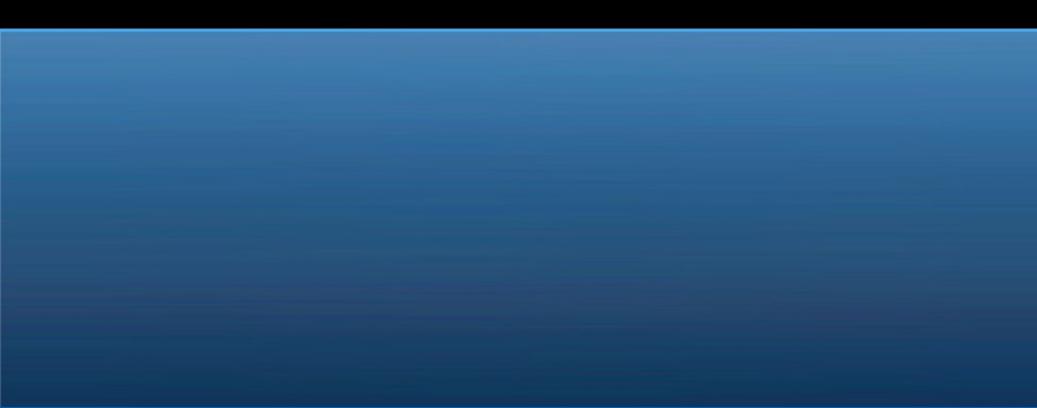
Managed via cgroup filesystem:

mkdir /dev/cgroup mount -t cgroup none /dev/cgroup mkdir /dev/cgroup/0 cd /dev/cgroup/0 echo \$\$ > tasks cat /proc/self/cgroup /etc/init.d/httpd start

Control Groups: history

- A feature known as cpusets was developed by big iron Bull/SGI guys
- Used to maintain process groups to NUMA nodes affinity
- Paul Menage generalized it
- Now cpusets is just one of the resource controllers

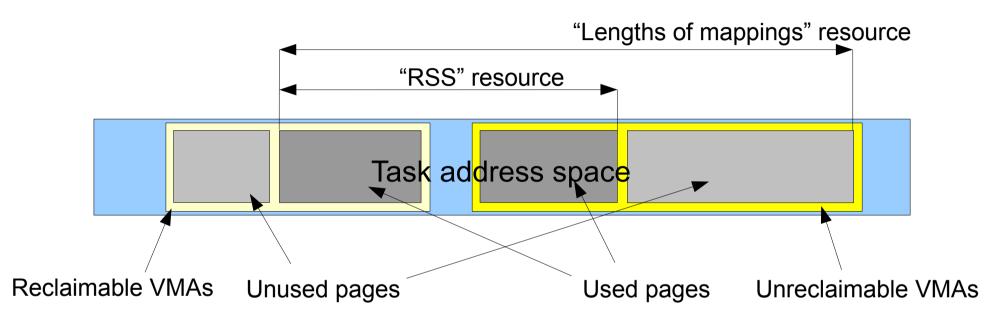
Memory Controller



Memory controller

- User memory:
 - RSS
 - Page cache
- Reclamation
 - Same as try_to_free_pages()
- OOM killer

User Memory



VMAs classification

- *unreclaimable*: private and anonymous
- reclaimable: shared file mappings

Pages classification

- *unused*: parts of mapped regions
- used: touched pages

MemCtrl: interface

```
# echo 4M > memory.limit_in_bytes
# cat memory.limit_in_bytes
4194304
# cat memory.usage_in_bytes
172032
# cat memory.max_usage_in_bytes
294912
# cat memory.failcnt
0
# cat memory.stat
```

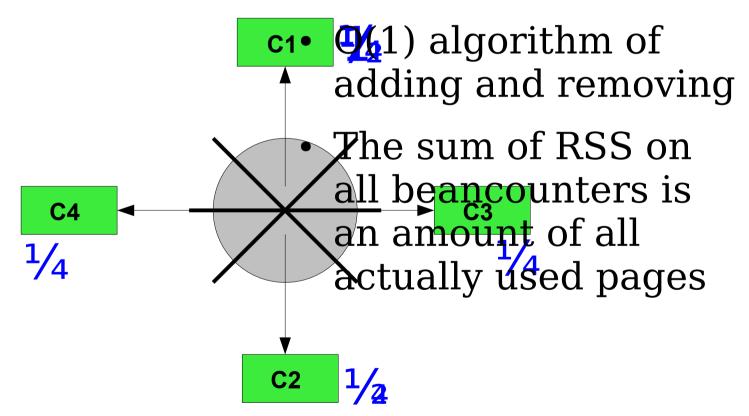
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Shared Pages accounting

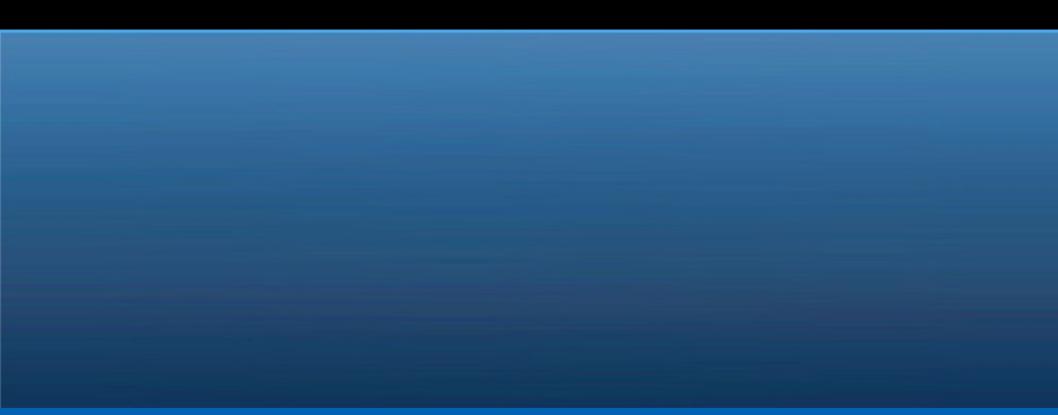
- Shared code/library segments
- Approaches:
 - Charge to the first user only (unfair)
 - Charge to all users (incorrect totals)
 - Charge a fraction to every user

Page fractions accounting

Algorithm benefits







Future a.k.a. TODO

- Shared pages accounting
- VMA (user mappings) length ctrl
- Kernel memory controller
- cgroups checkpoint/restart
- per-cgroup I/O priorities
- All that is available in OpenVZ; needs to be ported to mainstream

More Info

/usr/src/linux/Documentation/cgroups/*
/usr/src/linux/Documentation/controllers/*
 containers@linux-foundation.org

Questions? Comments?

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